

mum product and division errors occur as percentages of the operands, these algorithms are suited for high-speed hardware rather than GP computer applications.

One such application is real time digital filtering. The computations involved are usually sums of products of a variable and constant coefficients. Using the log-antilog algorithms gives complete freedom of coefficient selection. For single multiplications a cobweb array multiplier is simpler. However, for other configurations, such as a parallel digital filter bank, the log-antilog technique is less complex. Also, there are applications, such as multiplicative digital filters, where log and exponent conversions are necessary.

## REFERENCES

- [1] J. N. Mitchell, Jr., "Computer multiplication and division using binary logarithms," *IRE Trans. Electronic Computers*, vol. EC-11, pp. 512-517, August 1962.
- [2] M. Combet, H. Van Zonneveld, and L. Verbeek, "Computation of the base two logarithm of binary numbers," *IEEE Trans. Electronic Computers*, vol. EC-14, pp. 863-867, December 1965.
- [3] E. L. Hall, D. D. Lynch, and R. E. Young, "A digital modified discrete Fourier transform Doppler radar processor," *1968 EASCON Rec.*, pp. 150-159.
- [4] J. F. Kaiser and F. Kuo, *System Analysis by Digital Computer*. New York: Wiley, 1966, pp. 218-277.
- [5] C. M. Rader and B. Gold, "Digital filter design techniques in the frequency domain," *Proc. IEEE*, vol. 55, pp. 149-171, February 1967.
- [6] A. V. Oppenheim, R. W. Schaffer, and T. G. Stockham, Jr., "Nonlinear filtering of multiplied and convolved signals," *Proc. IEEE*, vol. 56, pp. 1264-1291, August 1968.

# A Generalization of the Fast Fourier Transform

J. A. GLASSMAN

**Abstract**—A procedure for factoring of the  $N \times N$  matrix representing the discrete Fourier transform is presented which does not produce shuffled data. Exactly one factor is produced for each factor of  $N$ , resulting in a fast Fourier transform valid for any  $N$ . The factoring algorithm enables the fast Fourier transform to be implemented in general with four nested loops, and with three loops if  $N$  is a power of two. No special logical organization, such as binary indexing, is required to unshuffle data. Included are two sample programs, one which writes the equations of the matrix factors employing the four key loops, and one which implements the algorithm in a fast Fourier transform for  $N$  a power of two. The algorithm is shown to be most efficient for  $N$  a power of two.

**Index Terms**—Cooley-Tukey algorithm, discrete Fourier transform, fast Fourier transform, mixed radix, spectral analysis.

THE fast Fourier transform, extensively covered in current journals [1]–[7] and popularly termed the Cooley-Tukey algorithm [4], can be generalized for any number of coefficients  $N$  by a factoring which does not shuffle the data. This factoring has three major effects on the fast Fourier transform. First, the transform is more easily explained, since the complexity of the tree graph to trace data is eliminated. Second, the mechanization is simplified since the final data need not be unshuffled, and third, the fast Fourier transform may be conveniently applied to any number of coefficients, although an application to anything but a power of two or four may be only of academic interest. The application of the fast Fourier transforms to any  $N$ , which may

be termed a mixed radix algorithm, has appeared in recent articles [7]–[9] but no other algorithm has been found which does not involve scrambled data. This paper develops the basic matrix, demonstrates the algorithm for its factorization, and illustrates the factoring and the resulting fast Fourier transform with programs for time-sharing operation.

## THE DISCRETE FOURIER TRANSFORM

The Fourier transform  $Y(w)$  of a function  $x(t)$  is defined by the relation

$$Y(w) = \int_{-\infty}^{\infty} x(t)e^{-j\omega t} dt. \quad (1)$$

If  $x(t)$  is sampled  $\mu$  times, a sampled function  $X^*(t)$  is produced, defined by

$$X^*(t) = \sum_{k=0}^{\mu-1} x(t)\delta(t - kT) \quad (2)$$

where  $\delta(t)$  is the Dirac delta function. The Fourier transform of  $X^*(t)$  is the discrete Fourier transform  $Y^*(w)$ :

$$\begin{aligned} Y^*(w) &= \int_{-\infty}^{\infty} x^*(t)e^{-j\omega t} dt \\ &= \int_{-\infty}^{\infty} \sum_{k=0}^{\mu-1} x(t)\delta(t - kT)e^{-j\omega t} dt \\ Y^*(w) &= \sum_{k=0}^{\mu-1} x(kT)e^{-j\omega kT}. \end{aligned} \quad (3)$$

If  $Y^*(w)$  is expressed at frequencies which are integer multiples  $n$  of  $1/NT$ , and if  $\mu$  is selected as some multiple  $m$  of  $N$ , the transform becomes

$$Y^*\left(\frac{n2\pi}{NT}\right) = \sum_{k=0}^{mN-1} x(kT) e^{-j(2\pi nk/N)} \quad (4)$$

which may be expressed as

$$Y^*\left(\frac{n2\pi}{NT}\right) = \sum_{k=0}^{N-1} \left( \sum_{i=0}^{m-1} x((k+iN)T) \right) e^{-j(2\pi nk/N)} \quad (5)$$

using the periodic property of the exponential function. For simplicity, the inner summation may be replaced by a derived function  $x'(t)$ , so that

$$Y^*\left(\frac{n2\pi}{NT}\right) = \sum_{k=0}^{N-1} x'(kT) e^{-j(2\pi nk/N)}. \quad (6)$$

The function of  $x'(t)$  is a folding of the function  $x(t)$ , implied by (5). If  $\mu = N$ ,  $x'(t)$  is identically  $x(t)$ . If  $x(t)$  is periodic with period  $NT$  and if it is sampled without noise, no advantage accrues to folding.

If  $N$  Fourier coefficients  $n=0$  to  $N-1$  are computed, the array of  $N$  coefficients may be expressed by

$$Y^* = M \cdot X \quad (7)$$

where

$$Y^* = \begin{bmatrix} Y^*(0) \\ Y^*(2\pi/NT) \\ Y^*(4\pi/NT) \\ \vdots \\ Y^*((N-1)2\pi/NT) \end{bmatrix}, \quad (8)$$

$$X = \begin{bmatrix} x'(0) \\ x'(T) \\ x'(2T) \\ \vdots \\ x'((N-1)T) \end{bmatrix}, \quad (9)$$

and  $M$  is the  $N \times N$  symmetric matrix which, with  $W = e^{-j2\pi/N}$ , is given by

$$M = \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & W^1 & W^2 & \cdots & W^{N-1} \\ 1 & W^2 & W^4 & \cdots & W^{N-2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & W^{N-1} & W^{N-2} & \cdots & W^1 \end{bmatrix} \quad (10)$$

for which the periodic property of the exponential has again been employed.

#### THE FACTORING ALGORITHM

If the Fourier transform is computed directly from (7),  $(N-1)^2$  complex multiplications are required. However, the  $M$  matrix may be factored producing one sparse matrix for each factor of  $N$ , which significantly reduces the number of complex multiplications required.

In general, if  $r$  is any factor of  $N$ , then letting  $N' = N/r$ , the  $j$ th row of the  $M$  matrix may be written as

$$m_j]^t = [m_{j1} m_{j2} \cdots m_{jr} W^{r(j-1)} \cdot m_{j1} W^{r(j-1)} \cdot m_{j2} \cdots W^{r(j-1)} \cdot m_{jr} \cdots W^{N'r(j-1)} \cdot m_{j1} \cdots W^{N'r(j-1)} \cdot m_{j2} \cdots W^{N'r(j-1)} \cdot m_{jr}] \quad (11)$$

where  $]$  denotes a column matrix and  $t$  denotes the transpose. To simplify the notation, let the partial row  $m_{j1} m_{j2} \cdots m_{jr}$ , equal  $\gamma_j$ ,

$$m_j]^t = [\gamma_j W^{r(j-1)} \cdot \gamma_j \cdots W^{N'r(j-1)} \cdot \gamma_j]. \quad (12)$$

The periodic property of  $W$  may be shown by writing

$$W^{r(j-1)} = W^{\{r(j-1)\}_{\text{MOD } N}}. \quad (13)$$

Since

$$\{r(j-1)\}_{\text{MOD } N} = r(j-1)_{\text{MOD } N'}, \quad (14)$$

(12) has  $N'$  distinct forms, one for each value of  $(j-1)_{\text{MOD } N'}$ . Thus,

$$m_j]^t = \begin{cases} [\gamma_j & \gamma_j & \cdots & \gamma_j] & , & (j-1)_{\text{MOD } N'} = 0 \\ [\gamma_j & W^r \cdot \gamma_j & \cdots & W^{r(N'-1)} \cdot \gamma_j] & , & (j-1)_{\text{MOD } N'} = 1 \\ [\gamma_j & W^{2r} \cdot \gamma_j & \cdots & W^{2r(N'-1)} \cdot \gamma_j] & , & (j-1)_{\text{MOD } N'} = 2 \\ \vdots & & & & & \\ [\gamma_j & W^{(N'-1)r} \cdot \gamma_j & \cdots & W^{(N'-1)r(N'-1)} \cdot \gamma_j] & , & (j-1)_{\text{MOD } N'} = N' - 1. \end{cases} \quad (15)$$

The  $N'$  relations of (15) yield to an obvious representation as the product of two matrices. Letting  $\phi_r$  be the row matrix of  $r$  zeros, the first  $N'$  values of  $j$  in (15) are equivalent to

$$\begin{bmatrix} \gamma_1 & \phi_r & \phi_r & \cdots & \phi_r \\ \phi_r & \gamma_2 & \phi_r & \cdots & \phi_r \\ \phi_r & \phi_r & \gamma_3 & \cdots & \phi_r \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \phi_r & \phi_r & \phi_r & \cdots & \gamma_{N'} \end{bmatrix} \begin{bmatrix} U_r & U_r & U_r & \cdots & U_r \\ U_r & W^{1r} \cdot U_r & W^{2r} \cdot U_r & \cdots & W^{(N'-1)r} \cdot U_r \\ U_r & W^{2r} \cdot U_r & W^{4r} \cdot U_r & \cdots & W^{(N'-2)r} \cdot U_r \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ U_r & W^{(N'-1)r} \cdot U_r & W^{(N'-2)r} \cdot U_r & \cdots & W^r \cdot U_r \end{bmatrix} \quad (16)$$

where  $U_r$  is the  $r \times r$  unit matrix. The first factor is an  $N' \times N'$  array of  $1 \times r$  submatrices.

The  $M$  matrix is an  $r$ -fold repetition of (15). Additional rows can be added to the first factor of (16) until each row of the  $M$  matrix is formed in place. The  $N \times N'$  array of  $1 \times r$  submatrices produced in this manner is the first factor of the algorithm termed  $F_1$ .

$$F_1 = \begin{bmatrix} 1 & 1 & \dots & 1 & 0 & 0 & \dots & 0 & 0 & 0 & \dots & 0 & \dots & 0 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 & 1 & W^1 & \dots & W^{r-1} & 0 & 0 & \dots & 0 & \dots & 0 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 & 0 & 0 & \dots & 0 & 1 & W^2 & \dots & W^{2(r-1)} & \dots & 0 & 0 & \dots & 0 \\ & & & & & & & & \vdots & & & & & & & & \\ 0 & 0 & \dots & 0 & 0 & 0 & \dots & 0 & 0 & 0 & \dots & 0 & \dots & 1 & W^{N'-1} & \dots & W^{(N'-1)(r-1)} \\ 1 & W^{N'} & \dots & W^{N'(r-1)} & 0 & 0 & \dots & 0 & 0 & 0 & \dots & 0 & \dots & 0 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 & 1 & W^{N'+1} & \dots & W^{(N'+1)(r-1)} & 0 & 0 & \dots & 0 & \dots & 0 & 0 & \dots & 0 \\ & & & & & & & & \vdots & & & & & & & & \\ 0 & 0 & \dots & 0 & 0 & 0 & \dots & 0 & 0 & 0 & \dots & 0 & \dots & 1 & W^{N'r-1} & \dots & W^{(N'r-1)(r-1)} \end{bmatrix} \cdot (17)$$

Note that the first factor contains nonzero entries identical to the first  $r$  columns of the  $M$  matrix, but staggered in a cyclic fashion through the  $N \times N$  array.

The second factor of (16) is of the same form as the  $M$  matrix except that each entry is a matrix rather than a single element, and each entry involves a power of  $W^r$  rather than  $W$ . The derivation of the first factoring equally applies to an array of matrices, so that if  $N'$  is not prime, then the second factor of (16) may be similarly factored. Therefore the second factor is denoted by  $M'$  to emphasize the recursive relation of the factoring. Continuing the notation, the factoring is

$$M = F_1 \cdot M^1 \quad (18)$$

$$M = F_1 \cdot F_2 \cdot M'' \quad (19)$$

$$\vdots$$

$$M = F_1 \cdot F_2 \cdot F_3 \cdot \dots \cdot F_k \quad (20)$$

where  $k$  is the number of factors of  $N$ . Each factor  $F_i$  is formed by a cyclic staggering of the first  $r_i$  rows of submatrices from  $M^{i-1}$  through an  $N \times N$  array where  $r_i$  is the factor of  $N$  separated at the  $i$ th factoring.

If each factor is a prime factor of  $N$ , and the factoring is complete, a minimum number of multiplications results for this algorithm. For convenience, the procedure employed in this paper is to order the factors by magnitude, postmultiplying in decreasing order. The procedure may be verified by partitioning  $F_1$  and  $M'$  into  $r \times r$  submatrices and multiplying.

#### FACTORING EXAMPLES

Several sample factorings serve to illustrate the equivalence of the product of the factors to the original matrix. For  $N=6$ , the  $M$  matrix and its two factors are

$$M = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & W^1 & W^2 & W^3 & W^4 & W^5 \\ 1 & W^2 & W^4 & W^0 & W^2 & W^4 \\ \hline 1 & W^3 & W^0 & W^3 & W^0 & W^3 \\ 1 & W^4 & W^2 & W^0 & W^4 & W^2 \\ 1 & W^5 & W^4 & W^3 & W^2 & W^1 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & W^1 & W^2 \\ 1 & W^2 & W^4 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 1 & W^3 & W^0 \\ 1 & W^4 & W^2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & W^5 & W^4 \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 \\ \hline 1 & 0 & 0 & W^3 & 0 & 0 \\ 0 & 1 & 0 & 0 & W^3 & 0 \\ 0 & 0 & 1 & 0 & 0 & W^3 \end{bmatrix}$$

These matrices have been partitioned into  $3 \times 3$  ( $r \times r$ ) submatrices to emphasize the technique for verifying the factoring. Two factors are also found for  $N=15$ , as follows.

$$F_1 = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & W^1 & W^2 & W^3 & W^4 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & W^2 & W^4 & W^6 & W^8 \\ 1 & W^3 & W^6 & W^9 & W^{12} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & W^4 & W^8 & W^{12} & W^{15} & 0 & 0 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & W^5 & W^{10} & W^0 & W^5 \\ 1 & W^6 & W^{12} & W^3 & W^9 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & W^7 & W^{14} & W^6 & W^{13} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & W^8 & W^1 & W^9 & W^2 \\ 1 & W^9 & W^3 & W^{12} & W^6 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 & 0 & 1 & W^{10} & W^5 & W^0 & W^{10} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & W^{11} & W^7 & W^3 & W^{14} \\ 1 & W^{12} & W^9 & W^6 & W^3 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & W^{13} & W^{11} & W^9 & W^7 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & W^{14} & W^{13} & W^{12} & W^{11} \end{bmatrix}$$

$$F_2 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 \\ \hline 1 & 0 & 0 & 0 & 0 & W^5 & 0 & 0 & 0 & 0 & W^{10} & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & W^5 & 0 & 0 & 0 & 0 & W^{10} & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & W^5 & 0 & 0 & 0 & 0 & W^{10} & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & W^5 & 0 & 0 & 0 & 0 & W^{10} & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & W^5 & 0 & 0 & 0 & 0 & W^{10} \\ \hline 1 & 0 & 0 & 0 & 0 & W^{10} & 0 & 0 & 0 & 0 & W^5 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & W^{10} & 0 & 0 & 0 & 0 & W^5 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & W^{10} & 0 & 0 & 0 & 0 & W^5 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & W^{10} & 0 & 0 & 0 & 0 & W^5 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & W^{10} & 0 & 0 & 0 & 0 & W^5 \end{bmatrix}$$

For  $N = 16$ ,  $M = F_1 \cdot F_2 \cdot F_3 \cdot F_4$ . Omitting the zero entries and the partitioning, the four factors are derived as follows.

$$F_1 = \begin{bmatrix} 1 & & & & & & & & & & & & & & \\ & 1 & & & & & & & & & & & & & \\ & & 1 & W^1 & & & & & & & & & & & \\ & & & 1 & W^2 & & & & & & & & & & \\ & & & & 1 & W^3 & & & & & & & & & \\ & & & & & 1 & W^4 & & & & & & & & \\ & & & & & & 1 & W^5 & & & & & & & \\ & & & & & & & 1 & W^6 & & & & & & \\ & & & & & & & & 1 & W^7 & & & & & \\ 1 & W^8 & & & & & & & & & & & & & \\ & & 1 & W^9 & & & & & & & & & & & \\ & & & 1 & W^{10} & & & & & & & & & & \\ & & & & 1 & W^{11} & & & & & & & & & \\ & & & & & 1 & W^{12} & & & & & & & & \\ & & & & & & 1 & W^{13} & & & & & & & \\ & & & & & & & 1 & W^{14} & & & & & & \\ & & & & & & & & 1 & W^{15} \end{bmatrix}$$





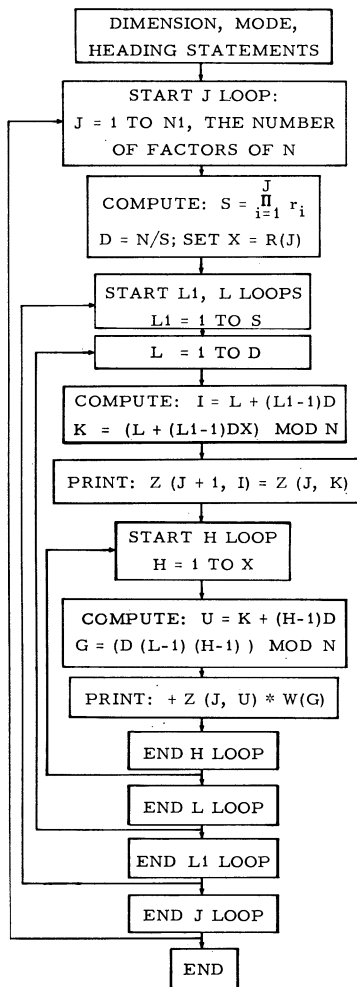


Fig. 1. Flow chart for the loops of the fast Fourier transform.

## LIST

FTLAMP 20:25 THURS. 07/11/68

```

100 DIMENSION Z(20,100),R(20)
110 INTEGER S,D,R,H,G,Z,U,X
120 4:READ,N,N1
130 PRINT,"EQUATIONS OF THE FAST FOURIER TRANSFORM,"
140 PRINT 2,N
150 2:FORMAT(11X,RHFOR N = ,12,3H = )
160 N2=N1-1
170 DO 1,J=1,N2
180 READ,R(J)
190 PRINT 1,R(J)
200 1:FORMAT(12,2H *)
210 READ,R(N1)
220 PRINT 3,R(N1)
230 3:FORMAT(12)
240 S=R(1)
250 DO 11,J=1,N1
260 X=R(J)
270 D=N/S
280 DO 31,L1=1,S
290 DO 31,L=1,D
300 61:FORMAT(2HZ(,12,1H,,12,6H) = Z(,12,1H,,12,2H) )
310 K=AMOD(L+(L1-1)*D*X,N)
320 I=L+(L1-1)*D
330 PRINT
340 PRINT 61,J+1,I,J,K,
350 DO 31,H=2,X
360 IF(AMOD(H-1,4)) 32,33,32
370 33:PRINT:PRINT
380 32:G=AMOD((L1-1)*D*(H-1),N)
390 U=X*(H-1)*D
400 PRINT 31,J,U,G,
410 31:FORMAT(4H+ Z(,12,1H,,12,4H)*W(,12,2H) )
420 IF(N1-1-J) 11
430 S=R(J+1)
440 11:PRINT+
450 PRINT++++GATA 4:END
460 $DATA
470 6,2,2,3,
480 8,3,2,2,2,
490 15,2,3,5,
500 16,4,2,2,2,2,
510 27,3,3,3,3,
520 30,3,2,3,5,
530 64,6,2,2,2,2,2,2
540 STOP

```

Fig. 2. Program to write fast Fourier transform equations.

duced by half. Thus for  $N = 2^\gamma$ , the required number of complex multiplications is

$$F(2^\gamma) = \frac{2^\gamma \gamma}{2} + 1 - 2^\gamma = 2^\gamma \left( \frac{\gamma}{2} - 1 \right) + 1. \quad (25)$$

A program has been written in FORTRAN for the GE-265 time-sharing system which writes the equations for the fast Fourier transform. The program illustrates the simple set of four loops necessary to mechanize the transform for any  $N$ , that is, employing a mixed radix, without special addressing logic. (For  $N$  a power of 2, the  $H$  loop becomes a single execution, and only three loops are required.) The flow chart for the basic routine is shown in Fig. 1; the program and sample results are shown in Figs. 2 and 3.

One of several programs written in FORTRAN employing the algorithm for radix 2 is listed in Fig. 4. This sample program stores the output complex coefficients in the input file and when the program is recalled, the inverse transform is taken as a verification. The output is the magnitude of the Fourier coefficients and is written into a file for computer plotting. A sample run for a

square pulse input is shown in Figs. 5 and 6 and the restored input function in Figs. 7 and 8. The program also includes folding of the input data as required by (5). Without loss of generality, the program outputs are computed for the regions of  $-NT/2 \leq t < NT/2$  and  $-1/2T \leq f < 1/2T$  to accommodate the symmetry available in the sample problem. The execution time for the time-sharing system appears to be dependent primarily upon the number of accesses to the disc for file data or for READ WRITE operations, and the program performs a transformation in approximately one minute for 128 complex coefficients and 128 input samples.

The fast Fourier transform shown in Fig. 4 employs a full buffer storage, the  $A()$  array, of length  $2N$ . This permits storing the full vector of complex numbers computed at each matrix multiplication, and may be a luxury in some applications. The basic algorithm requires buffer storage at each matrix multiplication, the extent depending upon the factor  $r$  being extracted and the position of the matrix in the sequence of matrix factors. The minimum buffer length has not been assessed; however, the maximum is clearly  $2N$  words.

FTL00P 20:27 THURS. 07/11/68

EQUATIONS OF THE FAST FOURIER TRANSFORM.  
FOR N = 6 = 2 \* 3

$$\begin{aligned} Z(2, 1) &= Z(1, 1) + Z(1, 4) * W(0) \\ Z(2, 2) &= Z(1, 2) + Z(1, 5) * W(0) \\ Z(2, 3) &= Z(1, 3) + Z(1, 6) * W(0) \\ Z(2, 4) &= Z(1, 1) + Z(1, 4) * W(3) \\ Z(2, 5) &= Z(1, 2) + Z(1, 5) * W(3) \\ Z(2, 6) &= Z(1, 3) + Z(1, 6) * W(3) \end{aligned}$$

$$\begin{aligned} Z(3, 1) &= Z(2, 1) + Z(2, 2) * W(0) + Z(2, 3) * W(0) \\ Z(3, 2) &= Z(2, 4) + Z(2, 5) * W(1) + Z(2, 6) * W(2) \\ Z(3, 3) &= Z(2, 1) + Z(2, 2) * W(2) + Z(2, 3) * W(4) \\ Z(3, 4) &= Z(2, 4) + Z(2, 5) * W(3) + Z(2, 6) * W(6) \\ Z(3, 5) &= Z(2, 1) + Z(2, 2) * W(4) + Z(2, 3) * W(8) \\ Z(3, 6) &= Z(2, 4) + Z(2, 5) * W(5) + Z(2, 6) * W(12) \end{aligned}$$

EQUATIONS OF THE FAST FOURIER TRANSFORM.  
FOR N = 8 = 2 \* 2 \* 2

$$\begin{aligned} Z(2, 1) &= Z(1, 1) + Z(1, 5) * W(0) \\ Z(2, 2) &= Z(1, 2) + Z(1, 6) * W(0) \\ Z(2, 3) &= Z(1, 3) + Z(1, 7) * W(0) \\ Z(2, 4) &= Z(1, 4) + Z(1, 8) * W(0) \\ Z(2, 5) &= Z(1, 1) + Z(1, 5) * W(4) \\ Z(2, 6) &= Z(1, 2) + Z(1, 6) * W(4) \\ Z(2, 7) &= Z(1, 3) + Z(1, 7) * W(4) \\ Z(2, 8) &= Z(1, 4) + Z(1, 8) * W(4) \end{aligned}$$

$$\begin{aligned} Z(3, 1) &= Z(2, 1) + Z(2, 3) * W(0) \\ Z(3, 2) &= Z(2, 2) + Z(2, 4) * W(0) \\ Z(3, 3) &= Z(2, 5) + Z(2, 7) * W(0) \\ Z(3, 4) &= Z(2, 6) + Z(2, 8) * W(0) \\ Z(3, 5) &= Z(2, 1) + Z(2, 3) * W(4) \\ Z(3, 6) &= Z(2, 2) + Z(2, 4) * W(4) \\ Z(3, 7) &= Z(2, 5) + Z(2, 7) * W(4) \\ Z(3, 8) &= Z(2, 6) + Z(2, 8) * W(4) \end{aligned}$$

$$\begin{aligned} Z(4, 1) &= Z(3, 1) + Z(3, 2) * W(0) \\ Z(4, 2) &= Z(3, 3) + Z(3, 4) * W(1) \\ Z(4, 3) &= Z(3, 5) + Z(3, 6) * W(2) \\ Z(4, 4) &= Z(3, 7) + Z(3, 8) * W(3) \\ Z(4, 5) &= Z(3, 1) + Z(3, 2) * W(4) \\ Z(4, 6) &= Z(3, 3) + Z(3, 4) * W(5) \\ Z(4, 7) &= Z(3, 5) + Z(3, 6) * W(6) \\ Z(4, 8) &= Z(3, 7) + Z(3, 8) * W(7) \end{aligned}$$

(a)

EQUATIONS OF THE FAST FOURIER TRANSFORM.  
FOR N = 15 = 3 \* 5

$$\begin{aligned} Z(2, 1) &= Z(1, 1) + Z(1, 6) * W(0) + Z(1, 11) * W(0) \\ Z(2, 2) &= Z(1, 2) + Z(1, 7) * W(0) + Z(1, 12) * W(0) \\ Z(2, 3) &= Z(1, 3) + Z(1, 8) * W(0) + Z(1, 13) * W(0) \\ Z(2, 4) &= Z(1, 4) + Z(1, 9) * W(0) + Z(1, 14) * W(0) \\ Z(2, 5) &= Z(1, 5) + Z(1, 10) * W(0) + Z(1, 15) * W(0) \\ Z(2, 6) &= Z(1, 1) + Z(1, 6) * W(5) + Z(1, 11) * W(10) \\ Z(2, 7) &= Z(1, 2) + Z(1, 7) * W(5) + Z(1, 12) * W(10) \\ Z(2, 8) &= Z(1, 3) + Z(1, 8) * W(5) + Z(1, 13) * W(10) \\ Z(2, 9) &= Z(1, 4) + Z(1, 9) * W(5) + Z(1, 14) * W(10) \\ Z(2, 10) &= Z(1, 5) + Z(1, 10) * W(5) + Z(1, 15) * W(10) \\ Z(2, 11) &= Z(1, 1) + Z(1, 6) * W(10) + Z(1, 11) * W(5) \\ Z(2, 12) &= Z(1, 2) + Z(1, 7) * W(10) + Z(1, 12) * W(5) \\ Z(2, 13) &= Z(1, 3) + Z(1, 8) * W(10) + Z(1, 13) * W(5) \\ Z(2, 14) &= Z(1, 4) + Z(1, 9) * W(10) + Z(1, 14) * W(5) \\ Z(2, 15) &= Z(1, 5) + Z(1, 10) * W(10) + Z(1, 15) * W(5) \end{aligned}$$

$$\begin{aligned} Z(3, 1) &= Z(2, 1) + Z(2, 2) * W(0) + Z(2, 3) * W(0) + Z(2, 4) * W(0) \\ &\quad + Z(2, 5) * W(0) \\ Z(3, 2) &= Z(2, 6) + Z(2, 7) * W(1) + Z(2, 8) * W(2) + Z(2, 9) * W(3) \\ &\quad + Z(2, 10) * W(4) \\ Z(3, 3) &= Z(2, 11) + Z(2, 12) * W(2) + Z(2, 13) * W(4) + Z(2, 14) * W(6) \\ &\quad + Z(2, 15) * W(8) \\ Z(3, 4) &= Z(2, 1) + Z(2, 2) * W(3) + Z(2, 3) * W(6) + Z(2, 4) * W(9) \\ &\quad + Z(2, 5) * W(12) \\ Z(3, 5) &= Z(2, 6) + Z(2, 7) * W(4) + Z(2, 8) * W(8) + Z(2, 9) * W(12) \\ &\quad + Z(2, 10) * W(16) \\ Z(3, 6) &= Z(2, 11) + Z(2, 12) * W(5) + Z(2, 13) * W(10) + Z(2, 14) * W(15) \\ &\quad + Z(2, 15) * W(20) \\ Z(3, 7) &= Z(2, 1) + Z(2, 2) * W(6) + Z(2, 3) * W(12) + Z(2, 4) * W(18) \\ &\quad + Z(2, 5) * W(24) \\ Z(3, 8) &= Z(2, 6) + Z(2, 7) * W(7) + Z(2, 8) * W(14) + Z(2, 9) * W(21) \\ &\quad + Z(2, 10) * W(28) \\ Z(3, 9) &= Z(2, 11) + Z(2, 12) * W(8) + Z(2, 13) * W(16) + Z(2, 14) * W(24) \\ &\quad + Z(2, 15) * W(32) \\ Z(3, 10) &= Z(2, 1) + Z(2, 2) * W(9) + Z(2, 3) * W(18) + Z(2, 4) * W(27) \\ &\quad + Z(2, 5) * W(36) \\ Z(3, 11) &= Z(2, 6) + Z(2, 7) * W(10) + Z(2, 8) * W(20) + Z(2, 9) * W(30) \\ &\quad + Z(2, 10) * W(40) \\ Z(3, 12) &= Z(2, 11) + Z(2, 12) * W(11) + Z(2, 13) * W(22) + Z(2, 14) * W(33) \\ &\quad + Z(2, 15) * W(44) \\ Z(3, 13) &= Z(2, 1) + Z(2, 2) * W(12) + Z(2, 3) * W(24) + Z(2, 4) * W(36) \\ &\quad + Z(2, 5) * W(48) \\ Z(3, 14) &= Z(2, 6) + Z(2, 7) * W(13) + Z(2, 8) * W(26) + Z(2, 9) * W(39) \\ &\quad + Z(2, 10) * W(52) \\ Z(3, 15) &= Z(2, 11) + Z(2, 12) * W(14) + Z(2, 13) * W(28) + Z(2, 14) * W(42) \\ &\quad + Z(2, 15) * W(56) \end{aligned}$$

(b)

EQUATIONS OF THE FAST FOURIER TRANSFORM.  
FOR N = 16 = 2 \* 2 \* 2 \* 2

$$\begin{aligned} Z(2, 1) &= Z(1, 1) + Z(1, 9) * W(0) \\ Z(2, 2) &= Z(1, 2) + Z(1, 10) * W(0) \\ Z(2, 3) &= Z(1, 3) + Z(1, 11) * W(0) \\ Z(2, 4) &= Z(1, 4) + Z(1, 12) * W(0) \\ Z(2, 5) &= Z(1, 5) + Z(1, 13) * W(0) \\ Z(2, 6) &= Z(1, 6) + Z(1, 14) * W(0) \\ Z(2, 7) &= Z(1, 7) + Z(1, 15) * W(0) \\ Z(2, 8) &= Z(1, 8) + Z(1, 16) * W(0) \\ Z(2, 9) &= Z(1, 1) + Z(1, 9) * W(8) \\ Z(2, 10) &= Z(1, 2) + Z(1, 10) * W(8) \\ Z(2, 11) &= Z(1, 3) + Z(1, 11) * W(8) \\ Z(2, 12) &= Z(1, 4) + Z(1, 12) * W(8) \\ Z(2, 13) &= Z(1, 5) + Z(1, 13) * W(8) \\ Z(2, 14) &= Z(1, 6) + Z(1, 14) * W(8) \\ Z(2, 15) &= Z(1, 7) + Z(1, 15) * W(8) \\ Z(2, 16) &= Z(1, 8) + Z(1, 16) * W(8) \end{aligned}$$

$$\begin{aligned} Z(3, 1) &= Z(2, 1) + Z(2, 5) * W(0) \\ Z(3, 2) &= Z(2, 2) + Z(2, 6) * W(0) \\ Z(3, 3) &= Z(2, 3) + Z(2, 7) * W(0) \\ Z(3, 4) &= Z(2, 4) + Z(2, 8) * W(0) \\ Z(3, 5) &= Z(2, 9) + Z(2, 13) * W(4) \\ Z(3, 6) &= Z(2, 10) + Z(2, 14) * W(4) \\ Z(3, 7) &= Z(2, 11) + Z(2, 15) * W(4) \\ Z(3, 8) &= Z(2, 12) + Z(2, 16) * W(4) \\ Z(3, 9) &= Z(2, 1) + Z(2, 5) * W(8) \\ Z(3, 10) &= Z(2, 2) + Z(2, 6) * W(8) \\ Z(3, 11) &= Z(2, 3) + Z(2, 7) * W(8) \\ Z(3, 12) &= Z(2, 4) + Z(2, 8) * W(8) \\ Z(3, 13) &= Z(2, 9) + Z(2, 13) * W(12) \\ Z(3, 14) &= Z(2, 10) + Z(2, 14) * W(12) \\ Z(3, 15) &= Z(2, 11) + Z(2, 15) * W(12) \\ Z(3, 16) &= Z(2, 12) + Z(2, 16) * W(12) \end{aligned}$$

$$\begin{aligned} Z(4, 1) &= Z(3, 1) + Z(3, 3) * W(0) \\ Z(4, 2) &= Z(3, 2) + Z(3, 4) * W(1) \\ Z(4, 3) &= Z(3, 5) + Z(3, 7) * W(2) \\ Z(4, 4) &= Z(3, 6) + Z(3, 8) * W(3) \\ Z(4, 5) &= Z(3, 9) + Z(3, 11) * W(4) \\ Z(4, 6) &= Z(3, 10) + Z(3, 12) * W(5) \\ Z(4, 7) &= Z(3, 13) + Z(3, 15) * W(6) \\ Z(4, 8) &= Z(3, 14) + Z(3, 16) * W(7) \\ Z(4, 9) &= Z(3, 1) + Z(3, 3) * W(8) \\ Z(4, 10) &= Z(3, 2) + Z(3, 4) * W(9) \\ Z(4, 11) &= Z(3, 5) + Z(3, 7) * W(10) \\ Z(4, 12) &= Z(3, 6) + Z(3, 8) * W(11) \\ Z(4, 13) &= Z(3, 9) + Z(3, 11) * W(12) \\ Z(4, 14) &= Z(3, 10) + Z(3, 12) * W(13) \\ Z(4, 15) &= Z(3, 13) + Z(3, 15) * W(14) \\ Z(4, 16) &= Z(3, 14) + Z(3, 16) * W(15) \end{aligned}$$

(c)

$$\begin{aligned} Z(5, 1) &= Z(4, 1) + Z(4, 2) * W(0) \\ Z(5, 2) &= Z(4, 3) + Z(4, 4) * W(1) \\ Z(5, 3) &= Z(4, 5) + Z(4, 6) * W(2) \\ Z(5, 4) &= Z(4, 7) + Z(4, 8) * W(3) \\ Z(5, 5) &= Z(4, 9) + Z(4, 10) * W(4) \\ Z(5, 6) &= Z(4, 11) + Z(4, 12) * W(5) \\ Z(5, 7) &= Z(4, 13) + Z(4, 14) * W(6) \\ Z(5, 8) &= Z(4, 15) + Z(4, 16) * W(7) \\ Z(5, 9) &= Z(4, 1) + Z(4, 2) * W(8) \\ Z(5, 10) &= Z(4, 3) + Z(4, 4) * W(9) \\ Z(5, 11) &= Z(4, 5) + Z(4, 6) * W(10) \\ Z(5, 12) &= Z(4, 7) + Z(4, 8) * W(11) \\ Z(5, 13) &= Z(4, 9) + Z(4, 10) * W(12) \\ Z(5, 14) &= Z(4, 11) + Z(4, 12) * W(13) \\ Z(5, 15) &= Z(4, 13) + Z(4, 14) * W(14) \\ Z(5, 16) &= Z(4, 15) + Z(4, 16) * W(15) \end{aligned}$$

EQUATIONS OF THE FAST FOURIER TRANSFORM.  
FOR N = 27 = 3 \* 3 \* 3

$$\begin{aligned} Z(2, 1) &= Z(1, 1) + Z(1, 10) * W(0) + Z(1, 19) * W(0) \\ Z(2, 2) &= Z(1, 2) + Z(1, 11) * W(0) + Z(1, 20) * W(0) \\ Z(2, 3) &= Z(1, 3) + Z(1, 12) * W(0) + Z(1, 21) * W(0) \\ Z(2, 4) &= Z(1, 4) + Z(1, 13) * W(0) + Z(1, 22) * W(0) \\ Z(2, 5) &= Z(1, 5) + Z(1, 14) * W(0) + Z(1, 23) * W(0) \\ Z(2, 6) &= Z(1, 6) + Z(1, 15) * W(0) + Z(1, 24) * W(0) \\ Z(2, 7) &= Z(1, 7) + Z(1, 16) * W(0) + Z(1, 25) * W(0) \\ Z(2, 8) &= Z(1, 8) + Z(1, 17) * W(0) + Z(1, 26) * W(0) \\ Z(2, 9) &= Z(1, 9) + Z(1, 18) * W(0) + Z(1, 27) * W(0) \\ Z(2, 10) &= Z(1, 1) + Z(1, 10) * W(9) + Z(1, 19) * W(18) \\ Z(2, 11) &= Z(1, 2) + Z(1, 11) * W(9) + Z(1, 20) * W(18) \\ Z(2, 12) &= Z(1, 3) + Z(1, 12) * W(9) + Z(1, 21) * W(18) \\ Z(2, 13) &= Z(1, 4) + Z(1, 13) * W(9) + Z(1, 22) * W(18) \\ Z(2, 14) &= Z(1, 5) + Z(1, 14) * W(9) + Z(1, 23) * W(18) \\ Z(2, 15) &= Z(1, 6) + Z(1, 15) * W(9) + Z(1, 24) * W(18) \\ Z(2, 16) &= Z(1, 7) + Z(1, 16) * W(9) + Z(1, 25) * W(18) \\ Z(2, 17) &= Z(1, 8) + Z(1, 17) * W(9) + Z(1, 26) * W(18) \\ Z(2, 18) &= Z(1, 9) + Z(1, 18) * W(9) + Z(1, 27) * W(18) \\ Z(2, 19) &= Z(1, 1) + Z(1, 10) * W(18) + Z(1, 19) * W(27) \\ Z(2, 20) &= Z(1, 2) + Z(1, 11) * W(18) + Z(1, 20) * W(27) \\ Z(2, 21) &= Z(1, 3) + Z(1, 12) * W(18) + Z(1, 21) * W(27) \\ Z(2, 22) &= Z(1, 4) + Z(1, 13) * W(18) + Z(1, 22) * W(27) \\ Z(2, 23) &= Z(1, 5) + Z(1, 14) * W(18) + Z(1, 23) * W(27) \\ Z(2, 24) &= Z(1, 6) + Z(1, 15) * W(18) + Z(1, 24) * W(27) \\ Z(2, 25) &= Z(1, 7) + Z(1, 16) * W(18) + Z(1, 25) * W(27) \\ Z(2, 26) &= Z(1, 8) + Z(1, 17) * W(18) + Z(1, 26) * W(27) \\ Z(2, 27) &= Z(1, 9) + Z(1, 18) * W(18) + Z(1, 27) * W(27) \end{aligned}$$

(d)

Fig. 3. Sample equations of the fast Fourier transform.



$$\begin{aligned}
Z(3,1) &= Z(2,1) + Z(2,4)*W(0) + Z(2,7)*W(0) \\
Z(3,2) &= Z(2,2) + Z(2,5)*W(0) + Z(2,8)*W(0) \\
Z(3,3) &= Z(2,3) + Z(2,6)*W(0) + Z(2,9)*W(0) \\
Z(3,4) &= Z(2,10) + Z(2,13)*W(0) + Z(2,16)*W(0) \\
Z(3,5) &= Z(2,11) + Z(2,14)*W(0) + Z(2,17)*W(0) \\
Z(3,6) &= Z(2,12) + Z(2,15)*W(0) + Z(2,18)*W(0) \\
Z(3,7) &= Z(2,19) + Z(2,22)*W(0) + Z(2,25)*W(0) \\
Z(3,8) &= Z(2,20) + Z(2,23)*W(0) + Z(2,26)*W(0) \\
Z(3,9) &= Z(2,21) + Z(2,24)*W(0) + Z(2,27)*W(0) \\
Z(3,10) &= Z(2,1) + Z(2,4)*W(9) + Z(2,7)*W(18) \\
Z(3,11) &= Z(2,2) + Z(2,5)*W(9) + Z(2,8)*W(18) \\
Z(3,12) &= Z(2,3) + Z(2,6)*W(9) + Z(2,9)*W(18) \\
Z(3,13) &= Z(2,10) + Z(2,13)*W(9) + Z(2,16)*W(18) \\
Z(3,14) &= Z(2,11) + Z(2,14)*W(9) + Z(2,17)*W(18) \\
Z(3,15) &= Z(2,12) + Z(2,15)*W(9) + Z(2,18)*W(18) \\
Z(3,16) &= Z(2,19) + Z(2,22)*W(9) + Z(2,25)*W(18) \\
Z(3,17) &= Z(2,20) + Z(2,23)*W(9) + Z(2,26)*W(18) \\
Z(3,18) &= Z(2,21) + Z(2,24)*W(9) + Z(2,27)*W(18) \\
Z(3,19) &= Z(2,1) + Z(2,4)*W(18) + Z(2,7)*W(9) \\
Z(3,20) &= Z(2,2) + Z(2,5)*W(18) + Z(2,8)*W(9) \\
Z(3,21) &= Z(2,3) + Z(2,6)*W(18) + Z(2,9)*W(9) \\
Z(3,22) &= Z(2,10) + Z(2,13)*W(18) + Z(2,16)*W(9) \\
Z(3,23) &= Z(2,11) + Z(2,14)*W(18) + Z(2,17)*W(9) \\
Z(3,24) &= Z(2,12) + Z(2,15)*W(18) + Z(2,18)*W(9) \\
Z(3,25) &= Z(2,19) + Z(2,22)*W(18) + Z(2,25)*W(9) \\
Z(3,26) &= Z(2,20) + Z(2,23)*W(18) + Z(2,26)*W(9) \\
Z(3,27) &= Z(2,21) + Z(2,24)*W(18) + Z(2,27)*W(9)
\end{aligned}$$

$$\begin{aligned}
Z(4,1) &= Z(3,1) + Z(3,2)*W(0) + Z(3,3)*W(0) \\
Z(4,2) &= Z(3,4) + Z(3,5)*W(0) + Z(3,6)*W(0) \\
Z(4,3) &= Z(3,7) + Z(3,8)*W(0) + Z(3,9)*W(0) \\
Z(4,4) &= Z(3,10) + Z(3,11)*W(0) + Z(3,12)*W(0) \\
Z(4,5) &= Z(3,13) + Z(3,14)*W(0) + Z(3,15)*W(0) \\
Z(4,6) &= Z(3,16) + Z(3,17)*W(0) + Z(3,18)*W(0) \\
Z(4,7) &= Z(3,19) + Z(3,20)*W(0) + Z(3,21)*W(0) \\
Z(4,8) &= Z(3,22) + Z(3,23)*W(0) + Z(3,24)*W(0) \\
Z(4,9) &= Z(3,25) + Z(3,26)*W(0) + Z(3,27)*W(0) \\
Z(4,10) &= Z(3,1) + Z(3,2)*W(9) + Z(3,3)*W(18) \\
Z(4,11) &= Z(3,4) + Z(3,5)*W(9) + Z(3,6)*W(18) \\
Z(4,12) &= Z(3,7) + Z(3,8)*W(9) + Z(3,9)*W(18) \\
Z(4,13) &= Z(3,10) + Z(3,11)*W(9) + Z(3,12)*W(18) \\
Z(4,14) &= Z(3,13) + Z(3,14)*W(9) + Z(3,15)*W(18) \\
Z(4,15) &= Z(3,16) + Z(3,17)*W(9) + Z(3,18)*W(18) \\
Z(4,16) &= Z(3,19) + Z(3,20)*W(9) + Z(3,21)*W(18) \\
Z(4,17) &= Z(3,22) + Z(3,23)*W(9) + Z(3,24)*W(18) \\
Z(4,18) &= Z(3,25) + Z(3,26)*W(9) + Z(3,27)*W(18) \\
Z(4,19) &= Z(3,1) + Z(3,2)*W(18) + Z(3,3)*W(9) \\
Z(4,20) &= Z(3,4) + Z(3,5)*W(18) + Z(3,6)*W(9) \\
Z(4,21) &= Z(3,7) + Z(3,8)*W(18) + Z(3,9)*W(9) \\
Z(4,22) &= Z(3,10) + Z(3,11)*W(18) + Z(3,12)*W(9) \\
Z(4,23) &= Z(3,13) + Z(3,14)*W(18) + Z(3,15)*W(9) \\
Z(4,24) &= Z(3,16) + Z(3,17)*W(18) + Z(3,18)*W(9) \\
Z(4,25) &= Z(3,19) + Z(3,20)*W(18) + Z(3,21)*W(9) \\
Z(4,26) &= Z(3,22) + Z(3,23)*W(18) + Z(3,24)*W(9) \\
Z(4,27) &= Z(3,25) + Z(3,26)*W(18) + Z(3,27)*W(9)
\end{aligned}$$

(e)

EQUATIONS OF THE FAST FOURIER TRANSFORM.  
FOR  $N = 30 = 2 * 3 * 5$ 

$$\begin{aligned}
Z(2,1) &= Z(1,1) + Z(1,16)*W(0) \\
Z(2,2) &= Z(1,2) + Z(1,17)*W(0) \\
Z(2,3) &= Z(1,3) + Z(1,18)*W(0) \\
Z(2,4) &= Z(1,4) + Z(1,19)*W(0) \\
Z(2,5) &= Z(1,5) + Z(1,20)*W(0) \\
Z(2,6) &= Z(1,6) + Z(1,21)*W(0) \\
Z(2,7) &= Z(1,7) + Z(1,22)*W(0) \\
Z(2,8) &= Z(1,8) + Z(1,23)*W(0) \\
Z(2,9) &= Z(1,9) + Z(1,24)*W(0) \\
Z(2,10) &= Z(1,10) + Z(1,25)*W(0) \\
Z(2,11) &= Z(1,11) + Z(1,26)*W(0) \\
Z(2,12) &= Z(1,12) + Z(1,27)*W(0) \\
Z(2,13) &= Z(1,13) + Z(1,28)*W(0) \\
Z(2,14) &= Z(1,14) + Z(1,29)*W(0) \\
Z(2,15) &= Z(1,15) + Z(1,30)*W(0) \\
Z(2,16) &= Z(1,1) + Z(1,16)*W(15) \\
Z(2,17) &= Z(1,2) + Z(1,17)*W(15) \\
Z(2,18) &= Z(1,3) + Z(1,18)*W(15) \\
Z(2,19) &= Z(1,4) + Z(1,19)*W(15) \\
Z(2,20) &= Z(1,5) + Z(1,20)*W(15) \\
Z(2,21) &= Z(1,6) + Z(1,21)*W(15) \\
Z(2,22) &= Z(1,7) + Z(1,22)*W(15) \\
Z(2,23) &= Z(1,8) + Z(1,23)*W(15) \\
Z(2,24) &= Z(1,9) + Z(1,24)*W(15) \\
Z(2,25) &= Z(1,10) + Z(1,25)*W(15) \\
Z(2,26) &= Z(1,11) + Z(1,26)*W(15) \\
Z(2,27) &= Z(1,12) + Z(1,27)*W(15) \\
Z(2,28) &= Z(1,13) + Z(1,28)*W(15) \\
Z(2,29) &= Z(1,14) + Z(1,29)*W(15) \\
Z(2,30) &= Z(1,15) + Z(1,30)*W(15)
\end{aligned}$$

$$\begin{aligned}
Z(3,1) &= Z(2,1) + Z(2,6)*W(0) + Z(2,11)*W(0) \\
Z(3,2) &= Z(2,2) + Z(2,7)*W(0) + Z(2,12)*W(0) \\
Z(3,3) &= Z(2,3) + Z(2,8)*W(0) + Z(2,13)*W(0) \\
Z(3,4) &= Z(2,4) + Z(2,9)*W(0) + Z(2,14)*W(0) \\
Z(3,5) &= Z(2,5) + Z(2,10)*W(0) + Z(2,15)*W(0) \\
Z(3,6) &= Z(2,16) + Z(2,21)*W(0) + Z(2,26)*W(0) \\
Z(3,7) &= Z(2,17) + Z(2,22)*W(0) + Z(2,27)*W(0) \\
Z(3,8) &= Z(2,18) + Z(2,23)*W(0) + Z(2,28)*W(0) \\
Z(3,9) &= Z(2,19) + Z(2,24)*W(0) + Z(2,29)*W(0) \\
Z(3,10) &= Z(2,20) + Z(2,25)*W(0) + Z(2,30)*W(0) \\
Z(3,11) &= Z(2,1) + Z(2,6)*W(9) + Z(2,11)*W(18) \\
Z(3,12) &= Z(2,2) + Z(2,7)*W(9) + Z(2,12)*W(18) \\
Z(3,13) &= Z(2,3) + Z(2,8)*W(9) + Z(2,13)*W(18) \\
Z(3,14) &= Z(2,4) + Z(2,9)*W(9) + Z(2,14)*W(18) \\
Z(3,15) &= Z(2,5) + Z(2,10)*W(9) + Z(2,15)*W(18) \\
Z(3,16) &= Z(2,16) + Z(2,21)*W(9) + Z(2,26)*W(18) \\
Z(3,17) &= Z(2,17) + Z(2,22)*W(9) + Z(2,27)*W(18) \\
Z(3,18) &= Z(2,18) + Z(2,23)*W(9) + Z(2,28)*W(18) \\
Z(3,19) &= Z(2,19) + Z(2,24)*W(9) + Z(2,29)*W(18) \\
Z(3,20) &= Z(2,20) + Z(2,25)*W(9) + Z(2,30)*W(18) \\
Z(3,21) &= Z(2,1) + Z(2,6)*W(18) + Z(2,11)*W(9) \\
Z(3,22) &= Z(2,2) + Z(2,7)*W(18) + Z(2,12)*W(9) \\
Z(3,23) &= Z(2,3) + Z(2,8)*W(18) + Z(2,13)*W(9) \\
Z(3,24) &= Z(2,4) + Z(2,9)*W(18) + Z(2,14)*W(9) \\
Z(3,25) &= Z(2,5) + Z(2,10)*W(18) + Z(2,15)*W(9) \\
Z(3,26) &= Z(2,16) + Z(2,21)*W(18) + Z(2,26)*W(9) \\
Z(3,27) &= Z(2,17) + Z(2,22)*W(18) + Z(2,27)*W(9) \\
Z(3,28) &= Z(2,18) + Z(2,23)*W(18) + Z(2,28)*W(9) \\
Z(3,29) &= Z(2,19) + Z(2,24)*W(18) + Z(2,29)*W(9) \\
Z(3,30) &= Z(2,20) + Z(2,25)*W(18) + Z(2,30)*W(9)
\end{aligned}$$

(f)

$$\begin{aligned}
Z(4,1) &= Z(3,1) + Z(3,2)*W(0) + Z(3,3)*W(0) + Z(3,4)*W(0) \\
Z(4,2) &= Z(3,5) + Z(3,6)*W(0) + Z(3,7)*W(0) + Z(3,8)*W(0) \\
Z(4,3) &= Z(3,9) + Z(3,10)*W(0) + Z(3,11)*W(0) + Z(3,12)*W(0) \\
Z(4,4) &= Z(3,13) + Z(3,14)*W(0) + Z(3,15)*W(0) + Z(3,16)*W(0) \\
Z(4,5) &= Z(3,17) + Z(3,18)*W(0) + Z(3,19)*W(0) + Z(3,20)*W(0) \\
Z(4,6) &= Z(3,21) + Z(3,22)*W(0) + Z(3,23)*W(0) + Z(3,24)*W(0) \\
Z(4,7) &= Z(3,25) + Z(3,26)*W(0) + Z(3,27)*W(0) + Z(3,28)*W(0) \\
Z(4,8) &= Z(3,29) + Z(3,30)*W(0) + Z(3,31)*W(0) + Z(3,32)*W(0) \\
Z(4,9) &= Z(3,1) + Z(3,2)*W(9) + Z(3,3)*W(18) + Z(3,4)*W(27) \\
Z(4,10) &= Z(3,5) + Z(3,6)*W(9) + Z(3,7)*W(18) + Z(3,8)*W(27) \\
Z(4,11) &= Z(3,9) + Z(3,10)*W(9) + Z(3,11)*W(18) + Z(3,12)*W(27) \\
Z(4,12) &= Z(3,13) + Z(3,14)*W(9) + Z(3,15)*W(18) + Z(3,16)*W(27) \\
Z(4,13) &= Z(3,17) + Z(3,18)*W(9) + Z(3,19)*W(18) + Z(3,20)*W(27) \\
Z(4,14) &= Z(3,21) + Z(3,22)*W(9) + Z(3,23)*W(18) + Z(3,24)*W(27) \\
Z(4,15) &= Z(3,25) + Z(3,26)*W(9) + Z(3,27)*W(18) + Z(3,28)*W(27) \\
Z(4,16) &= Z(3,29) + Z(3,30)*W(9) + Z(3,31)*W(18) + Z(3,32)*W(27) \\
Z(4,17) &= Z(3,1) + Z(3,2)*W(18) + Z(3,3)*W(9) + Z(3,4)*W(27) \\
Z(4,18) &= Z(3,5) + Z(3,6)*W(18) + Z(3,7)*W(9) + Z(3,8)*W(27) \\
Z(4,19) &= Z(3,9) + Z(3,10)*W(18) + Z(3,11)*W(9) + Z(3,12)*W(27) \\
Z(4,20) &= Z(3,13) + Z(3,14)*W(18) + Z(3,15)*W(9) + Z(3,16)*W(27) \\
Z(4,21) &= Z(3,17) + Z(3,18)*W(18) + Z(3,19)*W(9) + Z(3,20)*W(27) \\
Z(4,22) &= Z(3,21) + Z(3,22)*W(18) + Z(3,23)*W(9) + Z(3,24)*W(27) \\
Z(4,23) &= Z(3,25) + Z(3,26)*W(18) + Z(3,27)*W(9) + Z(3,28)*W(27) \\
Z(4,24) &= Z(3,29) + Z(3,30)*W(18) + Z(3,31)*W(9) + Z(3,32)*W(27) \\
Z(4,25) &= Z(3,1) + Z(3,2)*W(27) + Z(3,3)*W(18) + Z(3,4)*W(9) \\
Z(4,26) &= Z(3,5) + Z(3,6)*W(27) + Z(3,7)*W(18) + Z(3,8)*W(9) \\
Z(4,27) &= Z(3,9) + Z(3,10)*W(27) + Z(3,11)*W(18) + Z(3,12)*W(9) \\
Z(4,28) &= Z(3,13) + Z(3,14)*W(27) + Z(3,15)*W(18) + Z(3,16)*W(9) \\
Z(4,29) &= Z(3,17) + Z(3,18)*W(27) + Z(3,19)*W(18) + Z(3,20)*W(9) \\
Z(4,30) &= Z(3,21) + Z(3,22)*W(27) + Z(3,23)*W(18) + Z(3,24)*W(9) \\
Z(4,31) &= Z(3,25) + Z(3,26)*W(27) + Z(3,27)*W(18) + Z(3,28)*W(9) \\
Z(4,32) &= Z(3,29) + Z(3,30)*W(27) + Z(3,31)*W(18) + Z(3,32)*W(9)
\end{aligned}$$

(g)

LIST

FFT 20:19 THURS. 07/11/68

```

100 INTEGER GAMMA,SPACE,S,V,G,U,D,F1,F2,F3,F4,F5,F6
110 DIMENSION Z(512),A(512),W(256),Y(256)
120 $FILE DAPLNT,FTD0/FTD1/FTD2/FTD3/FTD4/FTD5/FTD6/FTD7/FTD8/FTD9
130 READ (2,11) N1,T1,SW1;IF(SW1)41,31,41
140 31:PRINT 32,N1,T1;33:42:41:PRINT 43,N1,T1;42:PRINT 51:PRINT 52,,
150 INPUT 4,SPACE;PHI=5.2831853/N;C=20/LOG(10);DELF=1/4/T1/SPACE
160 F1=M1/N/SPACE;F2=(M1+509)/510;F4=N/4+1
170 F5=(2*M1+509)/510;F6=(2*M1+5)/5;M=N/2;K1=510
180 GAMMA=1.93(4)/LOG(2)+1;M3=N*2;D=M/5;
190 IF(SW1) 71:D8 62;J=1;F2=IF(J-F5) 64;K1=M*SPACE*F1-(F2-1)*510
200 64:READ(2,11) (A(K),K=1,K1); D8 62;L=1;510
210 V=V+1;IF(V-M*SPACE*F1) 61,61;L=510;GOTO 62
220 61:IF(AMOD(V-1,SPACE)) 62,63,52:63:F3=AMOD(F3,M)+1
230 Z(2*F3-1)=Z(2*F3-1)+A(L);62;
240 33:42 72:71:00 72,J=1;F5=IF(J-F5) 77;K1=N*SPACE*F1*2-(F5-1)*510
250 77:READ(2,81) (A(K),K=1,K1);D8 72;L=1;510
260 V=V+1;IF(V-M*SPACE*F1) 75,76;L=510;GOTO 72;76:N2=AMOD(V-1,2*SPACE)
270 IF(N2) 74,73,74;73:F3=AMOD(F3,M)+1;Z(2*F3-1)=Z(2*F3-1)+A(L)
280 GOTO 72;74:IF(N2-1) 72,75,72;75:Z(2*F3)=Z(2*F3)+A(L); 72:
290 D8 91,J=1;W(J)=COS(PHI*(J-1))
300 91:W(M)=SIGN(SIN(PHI*(J-1)),.5-SW1)
310 D8 101,J=1,GAMMA
320 D8 111,L=1,S;D8 111,L=1,D
330 I=2*M+2*(L-1)*D;K=2*AMOD(L+(L-1)*D*2,N);U=K+2*M
340 G=(L-1)*D+1
350 B1=Z(U-1)*W(G)-Z(U)*W(G+M);B2=Z(U-1)*W(G+M)+Z(U)*W(G)
360 A(I-1)=Z(K-1)+B1;A(I)=Z(K)+B2
370 A(I+M-1)=Z(K-1)-B1;A(I+M)=Z(K)-B2
380 D=D/2;S=S*2
390 D8 101,J=1,N;101:Z(D)=A(J)
400 D8 141,J=1,N;141:Y(AMOD(J+M-1,N)+1)=SQRT(Z(2*J-1)*2+Z(2*J)*2)/F1/M
410 PRINT**PRINT 151:D8 194,J=1,N;194:V(J)=(J-M-1)*DELF;PRINT 152
420 PRINT**PRINT 161:J=1,N;161:IF(2*M-J) 171,173,171
430 171:PRINT 172,W(2*M-J),Y(2*M-J),C*(3*Y(2*M-J)-1),GOTO 162
440 173:PRINT 174,W(2*M-J),Y(2*M-J),162:IF(Y(2*M-J)) 191,192,191
450 191:PRINT 172,W(2*M-J),Y(2*M-J),C*LOG(Y(2*M-J));GOTO 161
460 192:PRINT 174,W(2*M-J),Y(2*M-J),161;
470 WRITE(1,192) 1;WRITE(1,193) N
480 WRITE(1,194) (W(J),J=1,N);WRITE(1,195) (Y(J),J=1,N)
490 WRITE(1,201);ENDFILE 1;REWIND 2;WRITE(2,11) N,DELF,1-SW1
500 WRITE(2,81) (Z(K),K=1,N3);ENDFILE 2
510 PRINT**PRINT 202,DELF
520 32:FORMAT(6HINPUT,15,22H REAL SAMPLES; PERIOD ,F4,6H UNITS)
530 43:FORMAT(6HINPUT,15,25H COMPLEX SAMPLES; PERIOD ,F4,6H UNITS)
540 51:FORMAT(68HHOW MANY COMPONENTS, N, ARE DESIRED AND WHAT SHOULD
550 +BE THE INTERVAL )
560 52:FORMAT(26H BETWEEN SAMPLES, SPACE)
570 151:FORMAT(X,7HSECONDS,3X,9HAMPLITUDE,4X,5HPPOWER,9X,
580 +8H SECONDS,3X,9HAMPLITUDE,4X,5HPPOWER)
590 152:FORMAT(7H OR CPS,6X,5HX1000,FX,2HDB,11X,6H8 CPS,6X,5HX1000,
600 +8X,2HDB)
610 172:FORMAT(FR,3X,3PF10,4,2X,6PF8,3X);81:FORMAT(6E11,5)
620 174:FORMAT(FR,3X,3PF10,4,5X,2H--1)
630 191:FORMAT(F10,41) 201:FORMAT(3F10,SPECTRUM . . . . .)
640 192:FORMAT(I1);193:FORMAT(I3)+1;FORMAT(I4,F4,11)
650 202:FORMAT(19H THE INTERVAL IS ,F10,6,7H UNITS.)
660 END

```

Fig. 4 Fast Fourier transform program for radix two.

FFT 19:35 THURS. 07/11/68

INPUT: 1280 REAL SAMPLES; PERIOD .0010 UNITS  
 HOW MANY COMPONENTS, N, ARE DESIRED AND WHAT SHOULD BE THE INTERVAL  
 BETWEEN SAMPLES, SPACE? 128,10

SECONDS OR CPS	AMPLITUDE X1000	POWER DB	SECONDS OR CPS	AMPLITUDE X1000	POWER DB
.000	62.5000	-24.082	.781	62.1054	-24.137
1.563	60.9304	-24.303	2.344	59.0011	-24.583
3.125	56.3603	-24.981	3.906	53.0661	-25.504
4.688	49.1909	-26.162	5.469	44.8193	-26.971
6.250	40.0456	-27.949	7.031	34.9719	-28.126
7.813	29.7053	-30.543	8.594	24.3553	-32.268
9.375	19.0305	-34.411	10.156	13.8369	-37.179
10.937	8.8745	-41.037	11.719	4.2350	-47.163
12.500	.0000	-202.378	13.281	3.7611	-48.494
14.062	6.9926	-43.107	14.844	9.6537	-40.306
15.625	11.7189	-38.622	16.406	13.1789	-37.692
17.187	14.0396	-37.053	17.969	14.3293	-36.880
18.750	14.0621	-37.039	19.531	13.3072	-37.518
20.312	12.1165	-38.332	21.094	10.5584	-39.528
21.875	8.7060	-41.202	22.656	6.6451	-43.550
23.437	4.4519	-47.009	24.219	2.2104	-53.111
25.000	.0000	-201.690	25.781	2.1045	-53.537
26.563	4.0350	-47.883	27.344	5.7321	-44.834
28.125	7.1464	-42.919	28.906	8.2399	-41.682
29.688	9.9862	-40.928	30.469	9.3720	-40.563
31.250	9.3960	-40.541	32.031	2.0695	-40.848
32.813	9.4150	-41.499	33.594	7.4658	-42.539
34.375	6.2639	-44.063	35.156	4.8592	-45.269
35.938	3.3072	-49.611	36.719	1.6672	-55.560
37.500	.0000	-210.033	38.281	1.6336	-55.737
39.063	3.1753	-49.954	39.844	4.5711	-46.800
40.625	5.7728	-44.772	41.406	6.7400	-43.427
42.187	7.4408	-42.568	42.969	7.8532	-42.099
43.750	7.9656	-41.976	44.531	7.7769	-42.184
45.313	7.2968	-42.737	46.094	6.5451	-43.682
46.875	5.5510	-45.113	47.656	4.3522	-47.226
48.438	2.9933	-50.477	49.219	1.5246	-56.337

THE INTERVAL IS .781250 UNITS.

AT LINE NO. 650: STOP END,RAN 692/6 SEC.

Fig. 5. Listing of square pulse transform.

SPECTRUM . . . . .

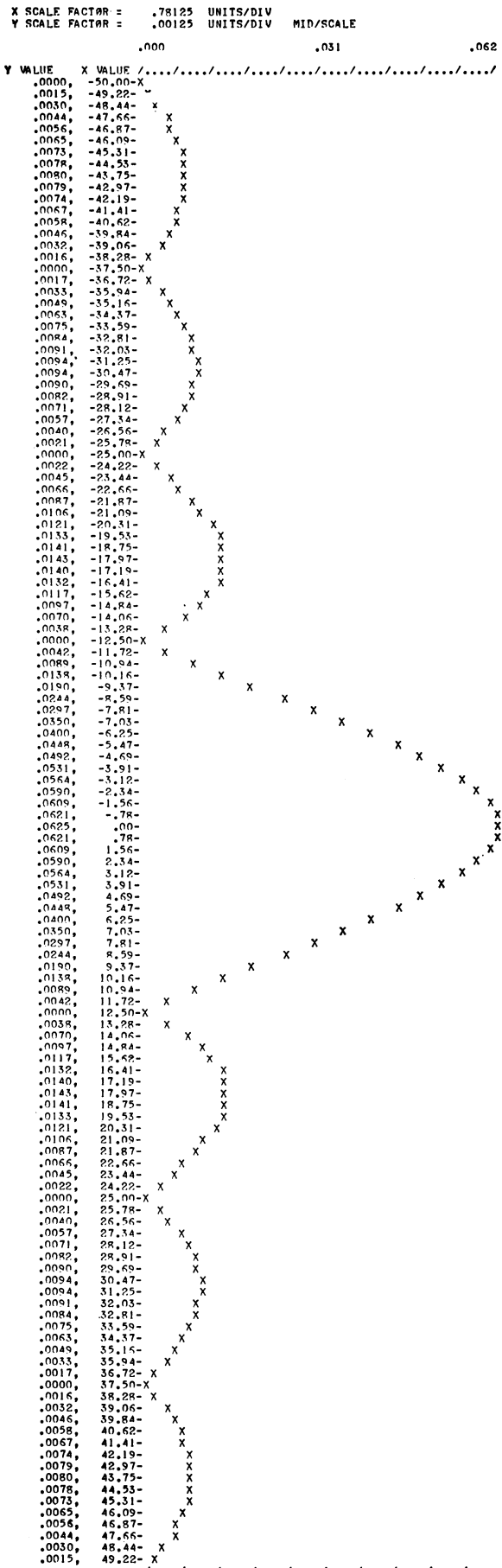


Fig. 6. Square pulse transform.

Fig. 7. Listing of restored pulse.

Fig. 8. Restored pulse.

## REFERENCES

- [1] W. T. Cochran, J. W. Cooley, D. L. Favin, H. D. Helms, R. A. Kaenel, W. W. Lang, G. C. Maling, Jr., D. E. Nelson, C. M. Rader, and P. D. Welch, "What is the fast Fourier transform," *Proc. IEEE*, vol. 55, pp. 1664-1674, October 1967.
- [2] V. W. Cooley, P. A. W. Lewis, and P. D. Welch, "Historical notes on the fast Fourier transform," *Proc. IEEE*, vol. 55, pp. 1675-1677, October 1967.
- [3] E. O. Brigham and R. W. Morrow, "The fast Fourier transform," *IEEE Spectrum*, vol. 4, pp. 63-70, December 1967.
- [4] V. W. Cooley and J. W. Tukey, "An algorithm for the machine calculation of complex Fourier series," *Math. Comp.*, vol. 19, pp. 297-301, 1965.
- [5] J. L. Shanks and T. W. Cairns, "Use of a digital convolution device to perform recursive filtering and the Cooley-Tukey algorithm," *IEEE Trans. Computers*, vol. C-17, pp. 943-949, October 1968.
- [6] R. C. Singleton, "ALGOL procedures for the fast Fourier transform," *Commun. ACM*, vol. 11, pp. 773-777, November 1968.
- [7] M. C. Pease, "An adaptation of the fast Fourier transform for parallel processing," *J. ACM*, vol. 15, pp. 252-264, April 1968.
- [8] R. C. Singleton, "An ALGOL procedure for the fast Fourier transform with arbitrary factors," *Commun. ACM*, vol. 11, pp. 777-779, November 1968.
- [9] W. M. Gentleman and G. Sande, "Fast Fourier transforms—for fun and profit," *1966 Fall Joint Computer Conf., AFIPS Proc.*, vol. 29. Washington, D. C.: Spartan, 1966, pp. 563-578.

# A Scheme for Synchronizing High-Speed Logic

## Part II

HERSCHEL H. LOOMIS, JR., MEMBER, IEEE

**Abstract**—In this paper we concern ourselves with the problem of obtaining high sequence rate sequential machines, machines which are constructed from realistic devices to operate at an input sequence rate which is independent of the machine complexity. To accomplish this result we have only to show a construction to realize acceptably synchronous devices from badly timed, restricted fan-in and fan-out devices. Once a complete set of synchronous devices is obtained, the results of Arden and Arthurs [1] apply and we know that any finite state machine has a realization using these devices which accepts input sequence members at a rate which is characteristic of the set of devices, not of the machine. The technique we propose for achieving this result is to produce a lattice of interconnected clock pulse sources called clock pulse propagators (CPPs). These devices generate clock pulses which are acceptably synchronized with respect to the outputs of neighboring CPPs but are not required to be in synchronization with some machine-wide standard as in current practice. Once it is established that such a network is possible, techniques already known can be applied in the utilization of the clock-pulses to synchronize logic and signals. Part I<sup>1</sup> of the paper concerns the analysis of CPP networks, and Part II covers the synthesis of sequential machines using CPP networks as clocking sources.

**Index Terms**—Badly timed devices, clocked logic, completeness of synchronous logic, high-speed logic, maximum rate construction.

### INTRODUCTION

IN this part of the paper, we examine the use of a lattice of clock sources in conjunction with badly timed logic and storage devices (latches) to produce

a complete set of synchronous logic devices for finite state machine synthesis purposes. The principal result of this part of the paper will be to provide a set of sufficient conditions for a set of badly timed logic devices to be complete in the sense of Loomis [1] and a design procedure for a clocked system. We will assume a lattice of clock pulse propagators (CPPs) as defined in Part I of this paper as our clock source, although the synchronizing method does not depend on the use of this source. Any source for which certain parameters can be bounded will suffice.

We first develop the constraints on the construction of the synchronous building blocks and on the physical spacing of the clock pulse lattice; then we present the theorem concerning completeness of badly timed devices. The paper is concluded with the presentation of a design procedure and an example.

### GENERAL DISCUSSION

The systems in which we shall use the CPP lattice as a clocking source shall consist of three elements:

- 1) the CPP lattice;
- 2) networks of gates which produce the desired switching functions;
- 3) latch circuits which synchronize the outputs of the gate networks with respect to the clock pulse sources.

Let us first consider the operation of the latch, and then we shall investigate the design constraints imposed by our maximum rate objective. The latch has been described in Maley and Earle [2]. It can be thought of

Manuscript received November 4, 1967; revised August 4, 1969. This paper is an extension of results reported in "A Theory of High-Speed Clocked Logic," presented at the 6th Annual Symposium on Switching Circuit Theory and Logical Design, October 1965. The work reported was supported by National Science Foundation Grant NSF GP3002.

H. H. Loomis, Jr., is with the Department of Electrical Engineering, University of California, Davis, Calif. 95616.

<sup>1</sup> H. H. Loomis, Jr., and M. R. McCoy, *IEEE Trans. Computers*, vol. C-19, pp. 39-47, January 1970.